

DARPA HACMS High Assurance Spiral



Spiral: Formal Approaches to Hardware & Software Design & Algorithm Verification

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Lecture based on joint work with CMU, UIUC, Drexel, and SpirlaGen, Inc.

The DARPA HACMS Program (K. Fisher)

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Source: DARPA-BAA-12-21 "High-Assurance Cyber Military Systems (HACMS)" Proposer's Day Slides by K. Fisher, HACMS Program Manager

The DARPA HACMS Program (K. Fisher)

DARPA Ubiquitous, Invisible, Networked, Computing Substrate

- In 2008, ~30 embedded processors per person in developed countries.
- In 2009, 98% of microprocessors were embedded [IEEE Computer `09]
- Trend: Networked embedded systems
- Vulnerabilities have economic and national security consequences. Extrapolating from *safety* failures:
 - June 10, 1999. Olympic Pipeline Company. 237K gallons of gasoline spilled. 3 deaths. >\$45M damages. [NTSB report]
 - Aug 14 2003. Northeast Blackout cost \$6B for 2 days of outage [DOE study]
 - April 26, 1986. Chernobyl Nuclear Disaster: >\$300B. Belarus alone: \$235B. [Chernobyl Forum]

August 14 2003 Northeast Blackout



Source: NOAA

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The growing connectivity between information systems, the Internet, and other infrastructures creates opportunities for attackers to disrupt telecommunications, electrical power, energy pipelines, refineries, financial networks, and other critical infrastructures."

-- Dennis C. Blair, Director of National Intelligence, *Annual Threat Assessment of the Intelligence Community for the Senate Select Committee on Intelligence, Statement for the Record,* February 12, 2009

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Our Approach: Model-Based High Assurance

Multi-sensor UGVs

- Multiple sensors: GPS, compass, accelerometer, IMU, etc.
- **Control:** waypoints, joystick vector
- Vehicle model: laws of physics, vehicle state
- Map data: Terrain, possible paths, obstacles

Assurance Through Consistency

- Model-based consistency checks
- Model vs. vehicle state
- Map-based path validation
- Exception signal if inconsistency threshold is exceeded







Assurance Through Consistency

- Model-based consistency checks Model vs. vehicle state
- Utilizes maps, physics, history, anticipated behavior, mission control
- Trusted virtual sensor output if model and sensors agree
- Exception if divergence beyond security threshold

High Assurance Controller



Assurance Through Guaranteed Controller Input and Output

- **Controller input:** virtual high-assurance sensor outputs
- Controller output: trusted or untrusted message to actuator
- **Controller algorithm:** PID or MPC, may use state, history and model
- **Failsafe:** use model-derived actuator setting if exception detected



Organization

Overview

Approach

- Example: Dynamic Window Monitor
- More HCOL examples
- Other research components
- Demos
- Concluding remarks





HCOL: Hybrid Control Operator Language

Sensor values and model-based predictions

Euler step: x^{t+h}

 $\mathbf{x}^{t+h} \approx \left[\, \mathbf{I}_{3} \, | h \, \mathbf{I}_{3} \, \right] (\mathbf{x}^{t} \oplus \mathbf{v}^{t+h})$

Numerical differentiation: v^{t+h}

 $\mathbf{v}^{t+h} \approx 1/h \Big[\mathbf{I}_3 | - \mathbf{I}_3 \Big] (\mathbf{x}^{t+h} \oplus \mathbf{x}^t)$

I₃: 3 x 3 identity matrix time step = matrix-vector product



Assurance through guaranteed controller input and output

- Declarative representation of physics, data and control algorithms
- Enables rule-based software synthesis and variant generation, verification and proof co-synthesis
- Extends Spiral's OL and SPL languages into the control domain



HCOL: Control Operator Examples

Time step residue: Disagreement between model and sensors

$$\mathbf{r}^{t+h} = \mathbf{R} \cdot (\mathbf{x}^t \oplus \mathbf{v}^t \oplus \mathbf{x}^{t+h} \oplus \mathbf{v}^{t+h})$$
$$\mathbf{R} = \begin{bmatrix} -\mathbf{I}_3 & h \mathbf{I}_3 & \mathbf{I}_3 & \mathbf{0}_3 \\ 1/h \mathbf{I}_3 & \mathbf{0}_3 & 1/h \mathbf{I}_3 & \mathbf{I}_3 \end{bmatrix}$$

Error operator: L₂ norm of time step residue $E_h : (\mathbf{s}^t, \mathbf{s}^{t+h}) \mapsto (\mathbf{s}^t \oplus \mathbf{s}^{t+h})^\top (\mathbf{R}^\top \mathbf{R}) (\mathbf{s}^t \oplus \mathbf{s}^{t+h})$ $\mathbf{s}^t = (\mathbf{x}^t \oplus \mathbf{v}^t), \quad \mathbf{s}^{t+h} = (\mathbf{x}^{t+h} \oplus \mathbf{v}^{t+h})$

 $\begin{aligned} & \mathsf{PID \ controller: \ Control \ velocity \ at \ set \ point \ \mathbf{v}_{0}} \\ & \begin{pmatrix} \mathbf{u}^{t} \\ \mathbf{s}^{t} \end{pmatrix} = \left(\begin{bmatrix} k_{p} \ \mathbf{I}_{e} \ | k_{i} \ \mathbf{I}_{3} \ | k_{d} / h \ \mathbf{I}_{3} \end{bmatrix} \begin{bmatrix} \mathbf{I}_{3} & \cdot & \cdot \\ \mathbf{I}_{3} & \cdot & \mathbf{I}_{3} \\ \mathbf{I}_{3} & -\mathbf{I}_{3} & \cdot \end{bmatrix} \begin{bmatrix} \mathbf{I}_{3} & -\mathbf{I}_{3} & \cdot \\ \cdot & \cdot & \mathbf{I}_{6} \end{bmatrix} \right) \oplus \begin{bmatrix} \mathbf{I}_{3} & -\mathbf{I}_{3} & \cdot \\ \mathbf{I}_{3} & -\mathbf{I}_{3} & \mathbf{I}_{3} \end{bmatrix} (\mathbf{v}_{0} \oplus \mathbf{v}^{t} \oplus \mathbf{s}^{t-h}) \\ & \mathbf{e}^{t} = \mathbf{v}^{0} - \mathbf{v}^{t}, \quad \mathbf{s}^{t} = \mathbf{e}^{t} \oplus \sum_{i=0}^{n-1} \mathbf{e}^{ih} \end{aligned}$

Usual PID controller definition: $\mathbf{u}^t = k_p \mathbf{e}^t + k_i \sum_{i=0}^{n-1} \mathbf{e}^{ih} + k_d \frac{\mathbf{e}^t - \mathbf{e}^{t-h}}{h}$

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Detection Through Feasible Region of State



Self-consistency equation

$$\mathcal{F}: \begin{pmatrix} \mathbf{x}^{t} \\ \mathbf{v}^{t} \\ \mathbf{x}^{t+h} \\ \mathbf{v}^{t+h} \end{pmatrix} \mapsto \begin{pmatrix} 1/h \left(\mathbf{x}^{t+h} - \mathbf{x}^{t} \right) \\ \mathbf{v}^{t+h} \end{pmatrix}$$

Inside a polyhedra

$$A_i \mathcal{F}(\vec{x}) - b_i \preceq \vec{0}$$

Test: attack-free, if $\mathcal{F}\left(\mathbf{s}^t \oplus \mathbf{s}^{t+h}\right) \in \bigcup_i \mathcal{P}_i$

Rule-Based Code Synthesis

High Level Rules: Transformations within high level abstraction

$$\begin{split} \mathbf{I}_{n} &\to \sum_{i=0}^{n-1} \mathbf{e}_{i}^{n} \mathbf{I}_{1}(\mathbf{e}_{i}^{n})^{\top} \\ \begin{bmatrix} \sum_{i=0}^{n-1} \mathbf{S}_{i} A_{i} \mathbf{G}_{i} \mid \sum_{i=0}^{n-1} \mathbf{S}_{i} B_{i} \mathbf{G}_{i} \end{bmatrix} &\to \sum_{i=0}^{n-1} \mathbf{S}_{i} \left(\begin{bmatrix} A_{i} \mid B_{i} \end{bmatrix} \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right) \mathbf{G}_{i} \\ \text{with } \mathbf{e}_{i}^{1 \times n} &= \begin{bmatrix} 0, \dots, 0, 1, 0, \dots, 0 \end{bmatrix} \end{split}$$

Code generation rules: Translate high level abstraction into code

$$\begin{aligned} \operatorname{Code}\left(y = (A B)x\right) &\to \left\{\operatorname{Decl}(t), \operatorname{Code}\left(t = Bx\right), \operatorname{Code}\left(y = At\right)\right\} \\ \operatorname{Code}\left(y = \left(\sum_{i=0}^{n-1} A_i\right)x\right) &\to \left\{y := \vec{0}, \operatorname{for}(i = 0..n - 1) \operatorname{Code}\left(y + A_ix\right)\right\} \\ \operatorname{Code}\left(y = e_i^{1 \times n}x\right) &\to y[0] := x[i] \\ \operatorname{Code}\left(y = e_i^{n \times 1}x\right) &\to \left\{y = \vec{0}, y[i] := x[0]\right\} \end{aligned}$$

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Co-Synthesis of Code and Correctness Proofs

Code generation: rule application until convergence

$\mathbf{y}^{t+h} = \left[\mathbf{I}_3 h \mathbf{I}_3\right] ($	$(\mathbf{x}^t \oplus \mathbf{v}^{t+h})$
	<pre>RuleSet := rec(SumSAG_In := Rule(@(I(@1)),(@, @1)->Let(i := Idx(@1), ISum(i, @1, e(@1, i) * I(1) * e(@1, i)^T))), SumDist :=,);</pre>
<pre>let(y:=var(TArray(T) func([inparam(xv) loop(i, [03] assign(nth(y)</pre>	Real, 3)), xv:=var(TArray(TReal, 6)), h := TReal(1/100),), outparam(y)], , chain(, i), add(nth(xv, i), mul(h, nth(xv, add(i,3)))))))))

Proof generation: trail of rule application

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Symbolic Rule Verification

 Rule replaces left-hand side by right-hand side when preconditions match

$$\mathbf{I}_n \to \sum_{i=0}^{n-1} \mathbf{e}_i^n \mathbf{I}_1(\mathbf{e}_i^n)^\top$$

Test rule by symbolically evaluating expressions before and after rule application and compare result

$$\mathbf{I}_{3} = \begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad = \mathbf{?} \quad \sum_{i=0}^{2} \mathbf{e}_{i}^{3} \mathbf{I}_{1} (\mathbf{e}_{i}^{3})^{\top} = \begin{bmatrix} 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \end{bmatrix}^{\top} + \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \quad = \mathbf{?} \quad \begin{bmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}^{\top} + \begin{bmatrix} 0 & 0 & 0 \end{bmatrix} \begin{bmatrix} 1 & 0 & 0 \end{bmatrix}^{\top} + \begin{bmatrix} 0 & 0 & 0 & 0 \end{bmatrix}^{\top}$$



Putting It All Together





Organization

- Overview
- Approach

Example: Dynamic Window Monitor

- More HCOL examples
- Other research components
- Demos
- Concluding remarks



Dynamic Window Safety Monitor



Dynamic Window Approach Primer







Algorithm Verified in KeYmaera

Theorem and proof

To Prove:
$$\psi_{ps} \to [dw_{ps}] \left((v_r = 0) \lor \left(\|p_r - p_o\| > \frac{v_r^2}{2b} + V \frac{v_r}{b} \right) \right)$$

 $\begin{aligned} dw_{ps} &\equiv (ctrl_{o} \mid\mid ctrl_{r}; dyn)^{*} \\ ctrl_{o} &\equiv v_{o} = (*, *); ? \mid\mid v_{o} \mid\mid \leq V \\ ctrl_{r} &\equiv (a_{r} := -b) \\ &\cup (?v_{r} = 0; a_{r} := 0; \omega_{r} := 0) \\ &\cup (a_{r} := *; ?-b \leq a_{r} \leq A; \omega_{r} := *; ?-\Omega \leq \omega_{r} \leq \Omega; \\ &p_{c} := (*, *); d_{r} := (*, *); p_{o} := (*, *); ?feasible \wedge safe) \end{aligned}$ $feasible &\equiv \mid\mid p_{r} - p_{c} \mid\mid > 0 \wedge \omega_{r} \mid\mid p_{r} - p_{c} \mid\mid = v_{r} \wedge d_{r} = \frac{(p_{r} - p_{c})^{\perp}}{\mid\mid p_{r} - p_{c} \mid\mid} \\ safe &\equiv \mid\mid p_{r} - p_{o} \mid\mid \infty > \frac{v_{r}^{2}}{2b} + V \frac{v_{r}}{b} + \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^{2} + \varepsilon(v_{r} + V)\right) \\ dyn &\equiv (t := 0; p_{r}^{x'} = v_{r}d_{r}^{x}, p_{r}^{y'} = v_{r}d_{r}^{y}, d_{r}^{x'} = -\omega_{r}d_{r}^{y}, d_{r}^{y'} = \omega_{r}d_{r}^{x}, \\ p_{o}^{x'} = v_{o}^{x}, p_{o}^{y'} = v_{o}^{y}, v_{r}' = a_{r}, \omega_{r}' = \frac{a_{r}}{\mid\mid p_{r} - p_{c} \mid\mid}, t' = 1 \\ &\& v_{r} \geq 0 \wedge t \leq \varepsilon) \end{aligned}$

Resulting safety monitor condition

$$\|p_r - p_o\|_{\infty} > \frac{v_r^2}{2b} + V\frac{v_r}{b} + \left(\frac{A}{b} + 1\right)\left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r + V)\right)$$





Proof/Code Co-Synthesis: HA Spiral



👷 High Assurance Spiral (Beta)	
Sigma-SPL expression:	_
Induction(3, Lambda([r1, r2], mul(r1, r2)), V(1.0))	
EXPANSION RULE: OLCompose	
SPL expression:	
ScalarProd(3, D) o Induction(3, Lambda([r1, r2], mul(r1, r2>), U(1.0>)	
Sigma-SPL expression:	
Reduction(3, (a, b) -> add(a, b), U(0.0), (arg) -> false) o PointWise(3, Lambda([r8, i2], mul(r8, nth(D, i2))) o Induction(3, Lambda([r1, r2], mul(r1, r2)), U(1.0))	
EXPANSION RULE: Reduction	
SPL expression:	
Reduction(2, (a, b) -> max(a, b), V(0.0), (arg) -> false)	
Sigma-SPL expression:	
Reduction(2, (a, b) -> max(a, b), V(0.0), (arg) -> false)	•

```
func(TInt, "transform", [ X, D ],
   decl([ u1, u2, u3, u4, u5, u6, u7, u8, w1, x1, x10, x1
      chain(
         ivenv(
            assign(u5, V([ V(0.0), V(0.0) ])),
            assign(u2, vcvt 64f32f(addsub 4x32f(vdup(Real
            assign(u1, V([ V(-1.0), V(1.0) ])),
            loop(i5, [ 0 .. 2 ],
               chain(
                  assign(x6, addsub 2x64f(vdup(add(RealEP
                  assign(x1, addsub 2x64f(V([ V(0.0), V(0
                  assign(x2, mul(x1, x6)),
                  assign(x3, mul(vushuffle 2x64f(x1, vpar
                  assign(x4, neg(min(x3, x2))),
                  assign(u3, add(max(vushuffle 2x64f(x4,
                  assign(u5, add(u5, u3)),
                  assign(x7, addsub 2x64f(V([ V(0.0), V(0
```

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Details: Formal Compilation

HCOL Breakdown Rules

SafeDist_{V,A,b,\varepsilon}(.,.,.) $\to (P[x, (a_0, a_1, a_2)](.) < d_{\infty}^2(.,.))(.,.,.)$

with
$$a_0 = \frac{1}{2b}, a_1 = \frac{V}{b} + \varepsilon \left(\frac{A}{b} + 1\right), a_2 = \left(\frac{A}{b} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon V\right)$$

 $d_{\infty}^{n}(.,.) \to \|.\|_{\infty}^{n} \circ (-)_{n}$

$$(\diamond)_n \rightarrow \mathsf{Pointwise}_{n \times n, (a,b) \mapsto a \diamond b}$$

 $\|.\|_{\infty}^{n} \to \operatorname{Reduction}_{n,(a,b)\mapsto\max(|a|,|b|)}$

 $<.,.>_n \rightarrow \mathsf{Reduction}_{n,(a,b)\mapsto a+b} \circ \mathsf{Pointwise}_{n \times n,(a,b)\mapsto ab}$

 $P[x, (a_0, \ldots, a_n)] \rightarrow < (a_0, \ldots, a_n), .> \circ (x^i)_n$

 $(x^i)_n \rightarrow \text{Induction}_{n,(a,b)\mapsto ab,1}$

Fully Expanded HCOL Expression

 $\begin{aligned} \mathsf{SafeDist}_{V,A,b,\varepsilon} \to \mathsf{Atomic}_{(x,y)\mapsto x < y} \\ & \circ \left(\Big(\mathsf{Reduction}_{3,(x,y)\mapsto x + y} \circ \mathsf{Pointwise}_{3,x\mapsto a_ix} \circ \mathsf{Induction}_{3,(a,b)\mapsto ab,1} \right) \\ & \times \Big(\mathsf{Reduction}_{2,(x,y)\mapsto \max(|x|,|y|)} \circ \mathsf{Pointwise}_{2\times 2,(x,y)\mapsto x - y} \Big) \Big) \end{aligned}$

```
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```



Final Synthesized C Code

```
int dwmonitor(float *X, double *D) {
     m128d u1, u2, u3, u4, u5, u6, u7, u8, x1, x10, x13, x14, x17, x18, x19, x2, x3, x4, x6, x7, x8, x9;
    int w1;
    ſ
        unsigned xm = mm getcsr();
        mm setcsr( xm & 0xffff0000 | 0x0000dfc0);
        u5 = mm set1 pd(0.0);
        u2 = mm cvtps pd( mm addsub ps( mm set1 ps(FLT MIN), mm set1 ps(X[0])));
        u1 = mm set pd(1.0, (-1.0));
        for(int i5 = 0; i5 <= 2; i5++) {
           x6 = mm addsub pd( mm set1 pd((DBL MIN + DBL MIN)), mm loaddup pd(&(D[i5])));
           x1 = _mm_addsub pd(mm set1 pd(0.0), u1);
           x^2 = mm mul pd(x^1, x^6);
           x3 = mm mul pd(mm shuffle pd(x1, x1, MM SHUFFLE2(0, 1)), x6);
            x4 = mm sub pd(mm set1 pd(0.0), mm min pd(x3, x2));
            u3 = mm add pd( mm max pd( mm shuffle pd(x4, x4, MM SHUFFLE2(0, 1)), mm max pd(x3, x2)), mm set1 pd(DBL MIN));
            u5 = mm add pd(u5, u3);
           x7 = mm addsub pd(mm set1 pd(0.0), u1);
           x8 = mm mul pd(x7, u2);
           x9 = mm mul pd(mm shuffle pd(x7, x7, MM SHUFFLE2(0, 1)), u2);
           x10 = mm sub pd(mm set1 pd(0.0), mm min pd(x9, x8));
            ul = mm add pd( mm max pd( mm shuffle pd(x10, x10, MM SHUFFLE2(0, 1)), mm max pd(x9, x8)), mm set1 pd(DBL MIN));
        }
       u6 = mm set1 pd(0.0);
        for(int i3 = 0; i3 <= 1; i3++) {
            u8 = mm cvtps pd( mm addsub ps( mm set1 ps(FLT MIN), mm set1 ps(X[(i3 + 1)])));
            u7 = _mm_cvtps_pd(_mm_addsub_ps(_mm_set1_ps(FLT_MIN), _mm_set1_ps(X[(3 + i3)])));
           x14 = mm add pd(u8, mm shuffle pd(u7, u7, MM SHUFFLE2(0, 1)));
           x13 = mm shuffle pd(x14, x14, MM SHUFFLE2(0, 1));
            u4 = mm shuffle pd( mm min pd(x14, x13), mm max pd(x14, x13), MM SHUFFLE2(1, 0));
            u6 = mm shuffle pd( mm min pd(u6, u4), mm max pd(u6, u4), MM SHUFFLE2(1, 0));
        }
        x17 = mm addsub pd(mm set1 pd(0.0), u6);
        x18 = mm addsub pd(mm set1 pd(0.0), u5);
        x19 = mm cmpge pd(x17, mm shuffle pd(x18, x18, MM SHUFFLE2(0, 1)));
        w1 = ( mm testc si128( mm castpd si128(x19), mm set epi32(0xffffffff, 0xffffffff, 0xffffffff, 0xffffffff)) -
             ( mm testnzc si128( mm castpd si128(x19), mm set epi32(0xffffffff, 0xffffffff, 0xffffffff, 0xffffffff, 0xffffffff))));
        asm nop;
       if (mm getcsr() & 0x0d) {
           mm setcsr( xm);
           return -1;
        }
        mm setcsr( xm);
    }
    return w1;
}
```

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Assembly Generated By Intel C Compiler

dwmonitor PROC	
sub	rsp, 120
vstmxcsr	DWORD PTR [112+rsp]
mov	r8d, DWORD PTR [112+rsp]
mov	eax, r8d
and	eax, -65536
or	eax, 57280
mov	DWORD PTR [112+rsp], eax
vldmxcsr	DWORD PTR [112+rsp]
vmovaps	<pre>xmm3, XMMWORD PTR [_2il0floatpacket.2]</pre>
vmovss	xmm0, DWORD PTR [rcx]
vshufps	<pre>xmm1, xmm0, xmm0, 0</pre>
vmovaps	<pre>xmm0, XMMWORD PTR [_2il0floatpacket.3]</pre>
vxorps	xmm5, xmm5, xmm5
vmovaps	xmm2, xmm5
vaddsubps	<pre>xmm4, xmm3, xmm1</pre>
vmovaps	<pre>xmm1, XMMWORD PTR [_2il0floatpacket.4]</pre>
vcvtps2pd	xmm4, xmm4
xor	eax, eax
vmovaps	XMMWORD PTR [32+rsp], xmm11
vmovaps	<pre>xmm11, XMMWORD PTR [_2il0floatpacket.5]</pre>
• • •	
vmovddup	<pre>xmm15, QWORD PTR [rdx+rax*8]</pre>
inc	rax
vaddsubpd	<pre>xmm13, xmm1, xmm15</pre>
vaddsubpd	<pre>xmm15, xmm5, xmm0</pre>
vminpd	<pre>xmm13, xmm14, xmm12</pre>
•••	
<100 more	lines>
• • •	
add	rsp, 120
ret	
ALIGN	16
dwmonitor ENDP	

64-bit mode AVX/VEX encoding 3 operand instructions SSE 4.1 1-1 mapping to C source 150 lines of assembly On SandyBridge: 100 – 240 cycles

30ns – 80ns @ 3 GHz



Spiral Interval Arithmetic Code Quality



SandyBridge CPU, Intel C Compiler, CompCert, APRON Interval Arithmetic Library



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Algorithms Formalized in HA Spiral

- Dynamic Window Approach Monitor
 Passive safety monitor, formally derived in KeYmaera
- Set calculus: Sensor self-consistency in state space
 Check that set of self-consistent true state values permitted
 by measurements is non-empty
- Multi-timescale Z-test for redundant sensors

Test for zero mean of difference between multiple sensors on multiple time scales

Mathematical infrastructure ROS code

Coordinate transformations, data filtering, ODE integration





Dynamic Window Safety Monitor



KeYmaera verification: monitors

Safety	Invariant + Safe Control	(RSS'13)
static	$\ p_r - p_o\ _{\infty} > rac{v_r^2}{2b} + \left(rac{A}{b} + 1 ight)\left(rac{A}{2}arepsilon^2 + arepsilon v_r ight)$	
passive	$v_r = 0 \vee \ p_r - p_o\ _{\infty} > \frac{v_r^2}{2b} + V\frac{v_r}{b} + \left(\frac{A}{b} + 1\right)\left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r)\right)$	(r + V)
+ sensor	$\ \hat{p}_r - p_o\ _{\infty} > \frac{v_r^2}{2b} + V\frac{v_r}{b} + \left(\frac{A}{b} + 1\right)\left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r + V)\right)$	$+ U_{\rho}$
+ disturb	$\ p_r - p_o\ _{\infty} > \frac{v_r^2}{2bU_m} + V \frac{v_r}{bU_m} + \left(\frac{A}{bU_m} + 1\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r + 1)\right) \left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r + 1)\right)$	V))
+ failure	$\ \hat{p}_r - p_o\ _{\infty} > \frac{v_r^2}{2b} + V\frac{v_r}{b} + \left(\frac{A}{b} + 1\right)\left(\frac{A}{2}\varepsilon^2 + \varepsilon(v+V)\right)$	$+ U_p + g\Delta$
friendly $ p_r - p_c $	$\ _{\infty} > \frac{v_r^2}{2b} + \frac{V^2}{2b_o} + V\left(\frac{v_r}{b} + \tau\right) + \left(\frac{A}{b} + 1\right)\left(\frac{A}{2}\varepsilon^2 + \varepsilon(v_r + V)\right)$)



Sensor Self-Consistency in State Space

Set calculus and approximation

Approximation through polytope



Inside a polytope \Rightarrow inside feasible set

$$A\vec{x} - b \preceq \vec{\mathsf{0}}$$

Time step and physics modeling

Intersect feasible sets of all sensors



Last intersection evolves with physics

$$p(t + \Delta) = p(t) + \Delta v(t) + \int_{t}^{t+\Delta} \int_{t}^{s} a(\tau) d\tau ds$$
$$v(t + \Delta) = v(t) + \int_{t}^{t+\Delta} a(s) ds$$



HCOL Specification and Expansion

HCOL Specification

$$\begin{aligned} \mathsf{ForAny}_{i=0}^{k-1} \left(\mathsf{InsidePoly}_{m,n}(\mathbf{A}_i,\mathbf{b}_i,.) \right) &: \mathbb{R}^n \to \mathbb{Z}_2 \\ x \mapsto (\exists i : \mathbf{A}_i x - \mathbf{b}_i \preceq \vec{\mathsf{0}}) \end{aligned}$$

Expansion into HCOL expression

$$\begin{aligned} & \operatorname{ForAny}_{i=0}^{k-1} \left(\operatorname{InsidePoly}_{m,n}(\mathbf{A}_{i},\mathbf{b}_{i},.) \right) \to \\ & \operatorname{Reduction}_{n,(a,b)\mapsto a\vee b} \\ & \circ \begin{bmatrix} \cdot \\ \cdot \end{bmatrix}_{i=0}^{k-1} \left(\operatorname{Reduction}_{n,(a,b)\mapsto a\wedge b} \\ & \circ \operatorname{Pointwise}_{n,x_{i}\mapsto x_{i}\leq 0} \circ \operatorname{Pointwise}_{n,x_{i}\mapsto x_{i}-b_{i}} \\ & \circ \begin{bmatrix} \cdot \\ \cdot \end{bmatrix}_{i=0}^{n-1} \left(\operatorname{Reduction}_{n,(a,b)\mapsto a+b} \circ \operatorname{Pointwise}_{n,x_{i}\mapsto a_{i}x_{i}} \right) \right) \end{aligned}$$





Multi-Timescale Z-Test

Receive a new residual value x. for Window size $w \in \{2^0, 2^1, ..., 2^{15}, \infty\}$ do Update number of samples $N_w \leftarrow N_w + 1$ Update the residual sample average \bar{x}_w to include x. if $N_w > w$ then Update \bar{x}_w to exclude the oldest sample in w. $N_w \leftarrow w$ end if {Compute a z-statistic for \bar{x}_w :} if $\bar{x}_w > \delta \mu_+$ then $z = \frac{\bar{x}_w - \mu_+}{\sigma/N_w}$ else if $\bar{x}_w < \delta \mu_-$ then $z = \frac{\bar{x}_w - \mu_-}{\sigma/N_w}$ else z = 0end if Extract p value using a Z-test on z. if $p < p_{thresh}$ then return Failure end if end for return Not Failure





HCOL Expansion

HCOL Operator Definition

 $Z_{w_{\max},\sigma,p_{\mathsf{thresh}},\delta,\mu_{+},\mu_{-}}^{n} : \mathbb{R}^{n} \to \mathbb{Z}_{2}$ $(x_{0},\ldots,x_{n-1}) \mapsto \bigvee_{w \in \{2^{0},2^{1},\ldots,w_{\max},n\}} \left(|z_{w}| > \Phi^{-1}\left(1 - \frac{p_{\mathsf{thresh}}}{2}\right)\right)$ with $N_{w} = \min(n,w)$ $\overline{x}_{w} = \sum_{i=\max(0,n-w)}^{n-1} x_{i}$ $z_{w} = \begin{cases} \sigma^{-1}N_{w}(\overline{x}_{w} - \mu_{+}) & \text{if } \overline{x}_{w} > \delta\mu_{+} \\ \sigma^{-1}N_{w}(\overline{x}_{w} - \mu_{-}) & \text{if } \overline{x}_{w} < \delta\mu_{-} \\ 0 & \text{else} \end{cases}$

HCOL Breakdown Rule

 $\begin{aligned} \mathsf{Z}^{n}_{w\max,\sigma,p_{\mathsf{thresh}},\delta,\mu_{+},\mu_{-}} &\to \mathsf{Reduction}_{2+\log_{2}w\max,(a,b)\mapsto a\lor b} \\ &\circ \mathsf{Pointwise}_{2+\log_{2}w\max,\tau^{p}\mathsf{thresh}_{\circ}z_{w}^{n,\sigma,\delta,\mu_{+},\mu_{-}}} \\ &\circ \begin{bmatrix} \cdot \\ \cdot \end{bmatrix}_{w\in\{2^{0},2^{1},\ldots,w\max,n\}} \overline{x}^{n}_{w} \end{aligned}$



Mathematial ROS Infrastructure Code

Example: (x,y) position from odometer

Euler step: x^{t+h}

 $\mathbf{x}^{t+h} \approx \left[\, \mathbf{I}_2 \, | h \, \mathbf{I}_2 \, \right] (\mathbf{x}^t \oplus \mathbf{v}^{t+h})$



Usual Euler definition:
$$(x^{t+h}, y^{t+h}) = (x^t + hv_x, y^t + hv_y)$$

PID controller: Control velocity at set point v₀

 $\begin{pmatrix} \mathbf{u}^t \\ \mathbf{s}^t \end{pmatrix} = \left(\begin{bmatrix} k_p \, \mathbf{I}_e \, | k_i \, \mathbf{I}_3 \, | k_d / h \, \mathbf{I}_3 \end{bmatrix} \begin{bmatrix} \mathbf{I}_3 & \cdot & \cdot \\ \mathbf{I}_3 & \cdot & \mathbf{I}_3 \\ \mathbf{I}_3 & -\mathbf{I}_3 & \cdot \end{bmatrix} \begin{bmatrix} \mathbf{I}_3 & -\mathbf{I}_3 & \cdot \\ \cdot & \cdot & \mathbf{I}_6 \end{bmatrix} \right) \oplus \begin{bmatrix} \mathbf{I}_3 & -\mathbf{I}_3 & \cdot \\ \mathbf{I}_3 & -\mathbf{I}_3 & \mathbf{I}_3 \end{bmatrix} (\mathbf{v}_0 \oplus \mathbf{v}^t \oplus \mathbf{s}^{t-h})$ $\mathbf{e}^t = \mathbf{v}^0 - \mathbf{v}^t, \quad \mathbf{s}^t = \mathbf{e}^t \oplus \sum_{i=0}^{n-1} \mathbf{e}^{ih}$

Usual PID controller definition: $\mathbf{u}^t = k_p \mathbf{e}^t + k_i \sum_{i=0}^{n-1} \mathbf{e}^{ih} + k_d \frac{\mathbf{e}^t - \mathbf{e}^{t-h}}{h}$



Electrical & Computer

High Assurance Spiral Code Generation

High Assurance Spiral (Beta)

```
- 🗆 ×
spiral> s := OLCompose(BinOp(3, Lambda([a, b], add(a, b))), PointWise(6, Lambda(
[x,i], cond(leg(i,Ū(2)), x, mul(x,h))));
BinOp(3, Lambda([ a1, b1 ], add(a1, b1))) o
PointWise(6, Lambda([ r1, i1 ], cond(leg(i1, U(2)), r1, mul(r1, param(TReal, "h"
DDDD
spiral> opts := HACMSopts.getOpts(rec(params := [h]));;
spiral> c := HACMSProof_Codegen(s, opts);;
spiral> c2 := Rewrite(Copy(c), RulesCodeHACMS, opts);
func(TVoid, "transform", [ Y, X, param(TReal, "h") ],
   dec1([ T2 ].
      chain(
         loop(i4, [ 0 .. 5 ],
            assign(nth(T2, i4), cond(leg(i4, U(2)), nth(X, i4), mul(nth(X, i4),
param(TReal, "h"))))
         ).
         loop(i5, [ 0 .. 2 ],
            assign(nth(Y, i5), add(nth(T2, i5), nth(T2, add(i5, U(3))))
         Э
      >
   )
spiral> c3 := Rewrite(Copy(c2), RulesCodeUnrollHACMS, opts);;
spiral> PrintCode("euler", c3, opts);
void euler(int *Y, double *X, double h) {
    double q10, q11, q12, q7, q8, q9;
    q7 = X[\bar{0}];
    a\bar{8} = X[1];
    \bar{q}9 = X[2];
    q10 = (X[3]*h);
    q11 = (X[4]*h);
    \alpha 12 = (X[5] + h);
    Y[0] = (q7 + q10);
    Y[1] = (a8 + a11);
    Y[2] = (\bar{q}9 + \bar{q}12);
spiral> _
```

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SpiralGen's High Assurance Spiral Tool Chain

<u>\$</u>	HACMS Build Progress
Sonx Sphinx	69 [Thread-3] INFO MetiTarskiLogger - MetiTarski ready 69 [Thread-3] INFO MetiTarskiLogger - Using axiom directory C:\metit-2.0tptp 71 [Thread-3] INFO MetiTarskiLogger - MetiTarski command arguments: [C:\metit-2.0\metit.exe,tptp, C:\metit 71 [Thread-3] INFO MetiTarskiLogger - Sending the following problem to MetiTarski:
KgY KeYmaera	<pre>% Auto-generated MetiTarski problem. % Number of variables: 18 fof(KeYmaera,conjecture, ![Y,A,DX,V,OY,DY,B,EP,OM,R] : ![DXUSCORE5DOLLARSK,OXUSCORE5DOLLARSK,RUSCORE5DOLLARSK,</pre>
Spiral	71 [Thread-3] ERROR MetiTarskiLogger - There was an Input/Output error while initialising the link with MetiT 72 [Thread-3] INFO MetiTarskiLogger - MetiTarski could not produce a proof! 77 [Thread-3] INFO MetiTarskiLogger - MetiTarski ready
Compile	78 [Thread-3] INFO MetiTarskiLogger - Using axiom directory C:\metit-2.0tptp 78 [Thread-3] INFO MetiTarskiLogger - MetiTarski command arguments: [C:\metit-2.0\metit.exe,tptp, C:\metit 79 [Thread-3] INFO MetiTarskiLogger - Sending the following problem to MetiTarski: % Auto-generated MetiTarski problem.
Deploy	<pre>% Number of variables: 18 fof(KeYmaera,conjecture, ![Y,A,DX,V,OY,DY,B,EP,OM,R] : ![DXUSCORE5DOLLARSK,OXUSCORE5DOLLARSK,RUSCORE5DOLLARSK, 79 [Thread-3] ERROR MetiTarskiLogger - There was an Input/Output error while initialising the link with MetiT</pre>
45%	<pre>79 [Thread-3] INFO MetiTarskiLogger - MetiTarski could not produce a proof! [DONE] KeYmaera script complete. Exiting KeYmaera Adding reference: eclipse.progress.monitor</pre>
	<pre>Adding reference: eclipse.progress.monitor - C:\Brian\SpiralGen\Repos\Beanstalk\Projects\HACMS\code_gen_tool\runtime-EclipseApplication/hsdf/src/dwmonit file C:/Brian/SpiralGen/Repos/Beanstalk/Projects/HACMS/code_gen_tool/runtime-EclipseApplication/hsdf/src/</pre>
	of DSP Algorithms http://www.spiral.net



Organization

- Overview
- Approach
- **Example: Dynamic Window Monitor**
- More HCOL examples
- Other research components
- Demos
- Concluding remarks



ModelPlex Runtime Validation

- ModelPlex ensures that proofs about models apply to real CPS
- Synthesize provably correct monitors to check CPS at runtime
- Correct-by-construction monitor conditions instead of manual annotation in models





Directional Collision Avoidance

Field of view and orientation

- Vehicle only responsible for collisions inside field of view
- Allows more aggressive driving: ignores obstacles outside visible area
- Narrow vision cone on straight lanes: fast with limited steering
- **Broad** vision cone at **intersections:** sharp turns at slow speed
- Multiple obstacle kinds
 - Pedestrians vs. other cars
 - Moveable vs. stationary
- Safety despite velocity uncertainty



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Formal Verification of PD Controller: Inverted Pendulum





Detection of Actuator + Sensor Attacks

 $x(k+1) = Ax(k) + Bu(k) + \Gamma e(k), \qquad \text{Actuator attacks}$ $y(k) = Cx(k) + \Psi e(k). \qquad \text{Sensor attacks}$

- Limitations of attack detection addressed as geometric control problems
- Detector performance depends on knowledge of system initial state
- One form of attack is undetectable when detector exactly knows system initial state:
 - Changes system's physical state (e.g. true velocity)
 - Does NOT change system sensor output (e.g., Odometer reading)





Estimating ABCar's Speed From Audio



- Recorded at HRL
- Multi microphone setup
- Audio classification
- Physics constraints (gear vs. speed)
- Good speed estimate (±2.5 km/h)







Execution Monitoring



Online execution monitoring to correct planning models about an adversary

Demonstrated in Robosoccer



Online Detection of Anomalies



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anom(R) = $\frac{P(obs|maximum likelihood anomaly in R)}{P(obs|no anomaly in R)}$



M. Veloso, J. P. Mendoza, and R. Simmons



Sensor Fusion And Data Consistency

Abnormal := not normal



Learn state-dependent bad data



Confidence in data



M. Veloso and J. P. Mendoza



Set-Bases Sensor Inconsistency Checks

- Fuse wheel encoders and GPS to detect inconsistencies
- Models noise and attack (strength, type)
- Matlab implementation calibrated with Carsim runs

Physics and noise model

$$p((k+1)\Delta) = p(k\Delta) + \int_{k\Delta}^{(k+1)\Delta} v(s)ds$$
$$y(k\Delta) = Hx(k\Delta) + v(k\Delta) + b(k\Delta)$$



If you have not seen CarSim or are a new user, view this video series to see how the software works.



Accelerating car with slight GPS attack





Electrical & Computer

Camera/Image Sensor Consistency

Projection and rotation matrices

M = KR[I, -C]

 $R = R_Y(Q_t + Q_{cv})R_X(Q_{ch})$

Cartoon and real image





Consistency check: compare cartoon image and camera image





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HACMS Phase 1 Demo on Landshark

- Setup: Drive Landshark with/without spoofing detection and obstacle avoidance, show impact of drive error and GPS attack
- Attack: Drift GPS to drive Landshark into obstacle while obstacle avoidance is engaged. Then show defense.
- **Tool:** Code synthesized with HA Spiral and KeYmaera/Sphinx
 - Run 1: no spoofing, no obstacle avoidance
 - Run 2: obstacle avoidance on
 - Run 3: obstacle avoidance,
 GPS spoofing attack
 - Run 4: obstacle avoidance + spoofing detection









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Calibrating The LandShark GPS









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Landshark Waypoint GPS Following

Con



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Summary: High Assurance Spiral

Problem and main idea

Co-synthesize high-quality code and proof for sensor-fusion based self-consistency algorithms



Results

- Four algorithms in HA Spiral formalized/in library dynamic window monitor, statistical tests, feasible state set test, infrastructure math code
- HA Spiral Tool/GUI ready for beta testers
- End-to-end proof/code co-synthesis and deployment deployed on Landshark and ABCar Simulator
- Rule based backend compiler proof of concept Spiral/Coq interface



int dwmonitor(float *X, double *D) m128d u1, u2, u3, u4, u5, u6, unsigned xm = mm getcsr(); mm setcsr(xm & 0xffff0000 | 0x u5 = mm set1 pd(0.0);u2 = mm cvtps pd(mm addsub ps mm set1 ps(FLT MIN), mm set1 p $\overline{u1} = mm \operatorname{set pd}(1.0, (-1.0));$ for (int i5 = 0; i5 <= 2; i5++) x6 = mm addsub_pd(_mm_set1 +DBL MIN)), mm loaddup x1 = mm addsub pd(mm set1 x2 = mm mul pd(x1, x6);asm nop; if (mm getcsr() & 0x0d) { mm setcsr(xm);

Approach



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